

Smallest snarks with oddness 4

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Snarks

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A **snark** is a connected cubic graph that has no **3**-edge-colouring.

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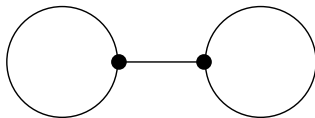
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- Four-Colour-Theorem/Problem
 - Cycle Double-Cover Conjecture
 - 5-Flow Conjecture
 - Fulkerson's Conjecture
- trivially true for 3-edge-colourable graphs
- open for snarks
- verified for graphs “close” to 3-edge-colourable graphs

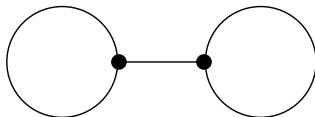
Small snarks

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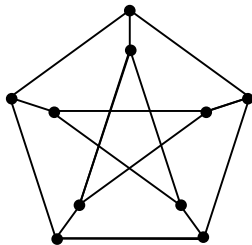


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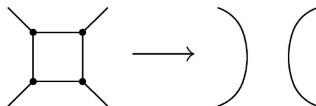
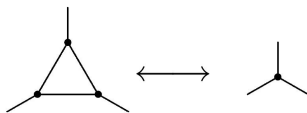
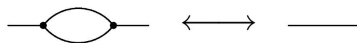
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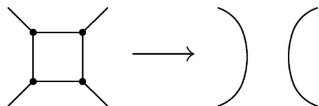
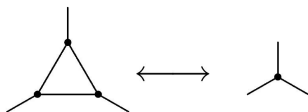
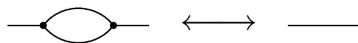
- smallest 2-connected snark



Non-trivial snarks

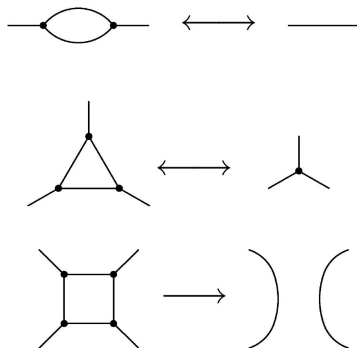


Non-trivial snarks



Similar simplifications for cycle-separating edge-cuts of size ≤ 3

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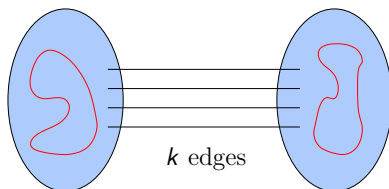
Similar simplifications for cycle-separating edge-cuts of size ≤ 3

\implies 'nontrivial' usually means

- cyclically 4-edge-connected, and
- girth > 4

Cyclic connectivity

Cyclic connectivity is the smallest number of edges whose removal leaves at least two components containing cycles.



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Theorem

The following statements are equivalent for a bridgeless cubic graph G :

- G is **3-edge-colourable**
- $\omega(G) = 0$
- $\rho(G) = 0$

Parity Lemma

Lemma (Parity lemma)

Let M be a k -pole edge-coloured with three colours. Let k_i be the number of dangling edges coloured with colour i . Then

$$k_1 \equiv k_2 \equiv k_3 \equiv k \pmod{2}.$$

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Corollary

For every snark G the following hold:

- $\rho(G) \geq 2$
- $\rho(G) = 2 \iff \omega(G) = 2$

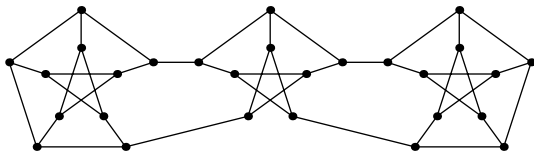
Small snarks with large oddness

Theorem

For every snark G there exists a snark G' of order not exceeding that of G such that $\omega(G') \geq \omega(G)$ and the girth of G' is at least 5.

Corollary

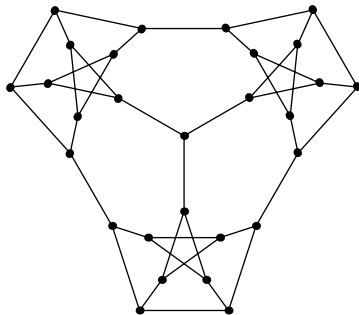
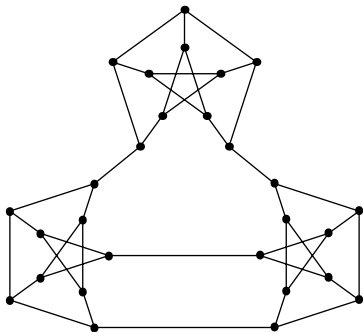
Every snark with oddness $\geq \omega$ and minimum order has girth at least 5.



Small snarks with $\omega \geq 4$

Theorem

The smallest order of a snark with $\omega \geq 4$ is 28. There are exactly three such snarks: two with connectivity 2 and one with connectivity 3.



Small **nontrivial** snarks with $\omega \geq 4$

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Brinkmann, Goedgebeur, Hägglund & Markström (2013)
generated all cyclically 4-edge-connected snarks up to 36 vertices.

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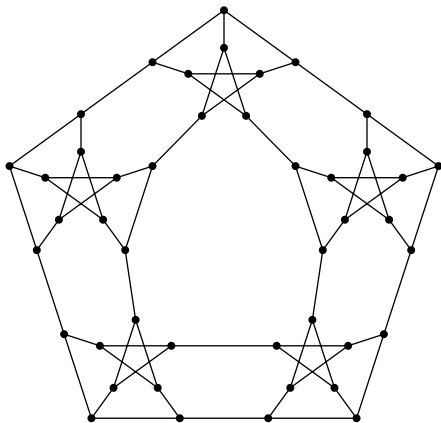
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Theorem (Brinkmann et al. (2013))

All cyclically 4-edge-connected snarks up to 36 vertices have $\omega = 2$.

Snark of order 44 and $\omega = 4$



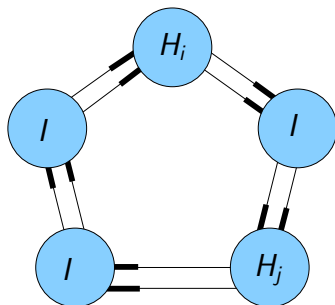
Constructed by Lukot'ka, Máčajová, Mazák & S. (2015)

Main result

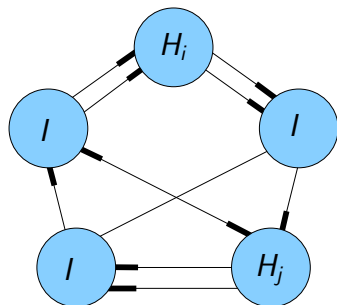
Theorem

- *The smallest number of vertices of a snark with cyclic connectivity 4 and $\omega \geq 4$ is 44.*
- *There are exactly 31 such snarks, all with $\omega = 4$.*

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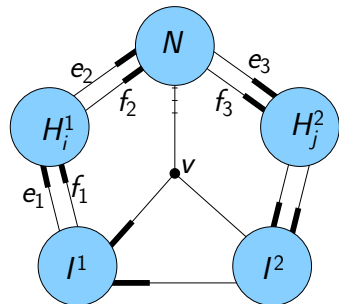


(a) Group 1a: 3 graphs

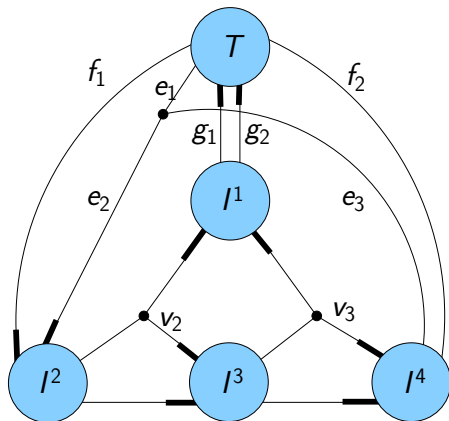


(b) Group 1b: 4 graphs

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(c) Group 4: 4 graphs



(d) Group 6: 2 graphs

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The proof has two main ingredients:

- Analysis of 3-edge-colourings conflicting on a cycle-separating 4-edge-cuts.
- A “closure theorem” for cyclically 4-edge-connected graphs of Andersen, Fleischner & Jackson (1988)

Closure Theorem

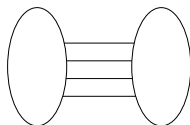
Theorem (Andersen, Fleischner & Jackson, 1988)

Let H be a cyclically 4-edge-connected cubic graph with a cycle-separating 4-edge-cut S . Then each component of $H - S$ can be extended to a cyclically 4-edge-connected graph by adding two adjacent vertices.

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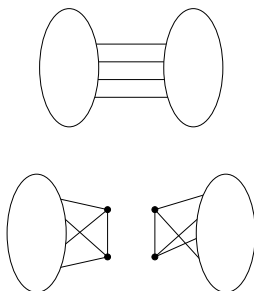
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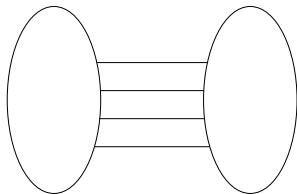
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Suppose to the contrary there exists a snark G on < 44 vertices with $\omega \geq 4$ and cyclic connectivity $= 4$.

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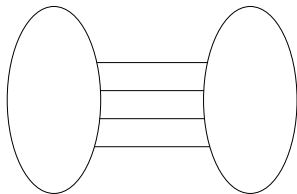
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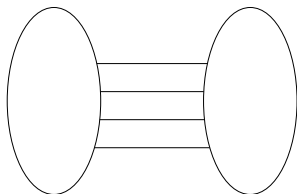
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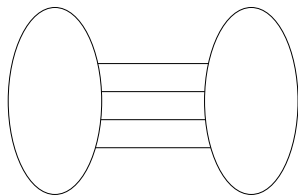
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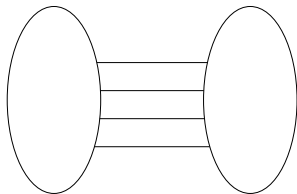
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- If both sides are colourable, then both have two types of colourings
 $\Rightarrow \rho(G) = 2 \Rightarrow \omega(G) = 2 \dots$ a contradiction.

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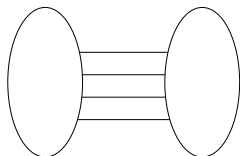


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- \Rightarrow At least one part is not 3-edge-colourable.

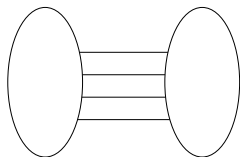
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We prove that each side can be extended to a snark by adding either

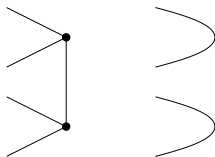
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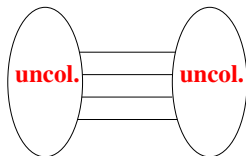


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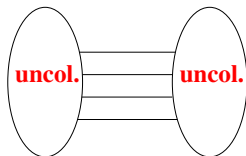


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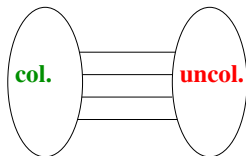
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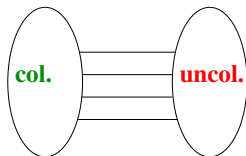
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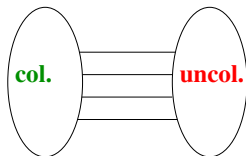
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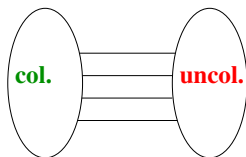
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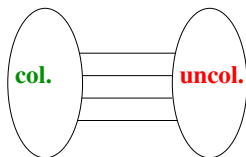
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 - ▶ By minimality, \tilde{G}_1 is also cyclically 4-edge-connected.

End of proof

- So far we have proved:

Every snark with $\omega \geq 4$, cyclic connectivity = 4, and minimum order can be decomposed into two smaller cyclically 4-edge-connected snarks.

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- If there is one on 42 or fewer vertices, it must arise by the reverse process from two snarks on at most 36 vertices.
- We have computationally verified all possible pairs of snarks and checked that none of them has $\omega \geq 4$.

Final remarks

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To prove a similar result for **nontrivial** snarks (rather than just those with cyclic connectivity = 4) we need to exclude the existence of cyclically 5-edge-connected snarks of orders 38 – 42 with $\omega \geq 4$.

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The existence of such snarks is unlikely.

Thank you!