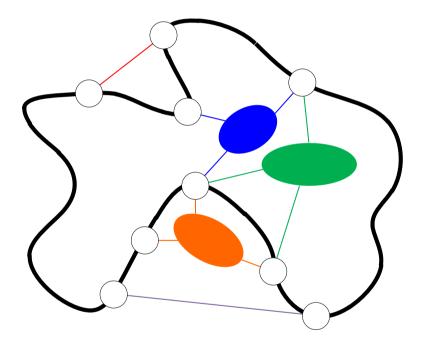
#### Longer Cycles in Essentially 4-Connected Planar Graphs



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joint work with Igor Fabrici, Jochen Harant and Samuel Mohr

## Circumference

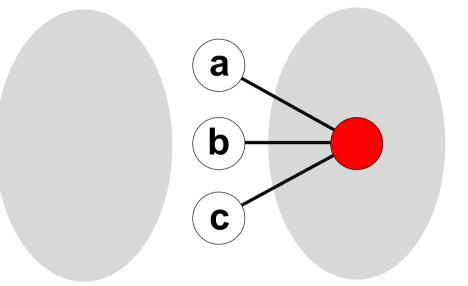
- G := polyhedral graph, i.e. planar and 3-connected
- circ(G) := length of a longest cycle of G (circumference)
  → If circ(G) = n := |V|, G has a Hamiltonian cycle.

What is the circumference of polyhedral graphs?

- Theorem [1963 Moon-Moser]: For infinitely many such graphs G, circ(G) ≤ 9n<sup>log<sub>3</sub> 2</sup> (log<sub>3</sub> 2 ≈ 0.631).
- Theorem [2002 Chen-Yu]: For all such G, there is a positive constant c such that circ(G) ≥ cn<sup>log<sub>3</sub> 2</sup>.
- Theorem [1956 Tutte]: Every 4-connected planar graph G is Hamiltonian.

## **Between 3- and 4-connectivity**

• A polyhedral graph G is essentially 4-connected if every 3-separator is the neighborhood of a single vertex.

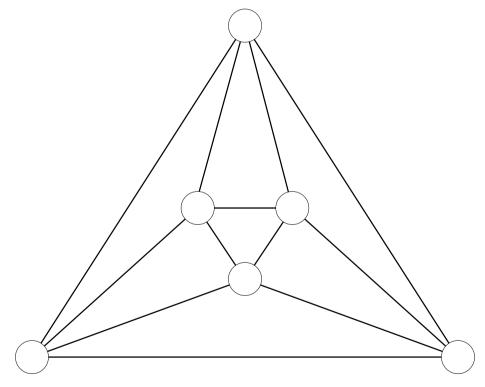


- Theorem [1992 Jackson-Wormald]: For every essentially 4-connected polyhedral graph G, circ(G) ≥ (2n+4)/5.
- Theorem [1976 Grünbaum-Malkevitch]: For every cubic essentially 4-connected polyhedral graph G, circ(G) ≥ 3n/4.

#### **But not Hamiltonian**

Theorem [2016 Fabrici-Harant-Jendrol', 1992 Jackson-Wormald]: For infinitely many (maximal planar) essentially 4-connected polyhedral graphs G,  $circ(G) \le 2(n+4)/3$ .

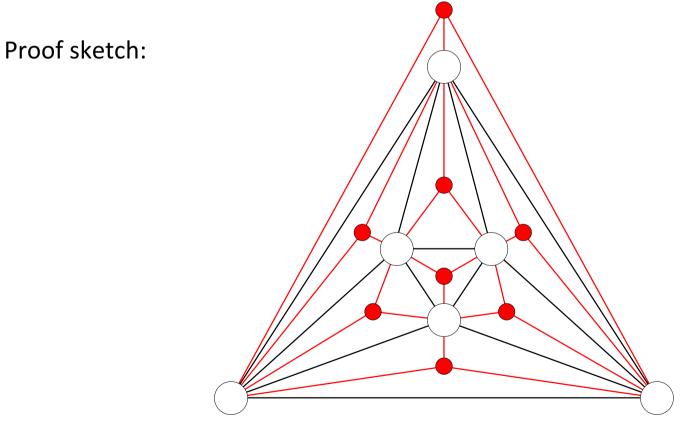
Proof sketch:



• Take any 4-connected maximal planar graph.

## **But not Hamiltonian**

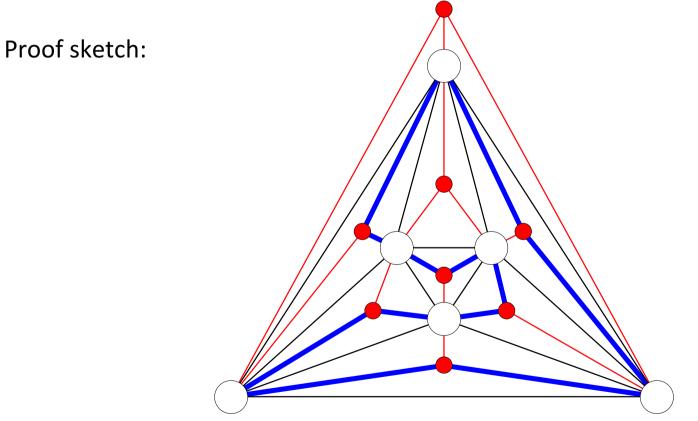
Theorem [2016 Fabrici-Harant-Jendrol', 1992 Jackson-Wormald]: For infinitely many (maximal planar) essentially 4-connected polyhedral graphs G,  $circ(G) \le 2(n+4)/3$ .



- Take any 4-connected maximal planar graph.
- Insert a vertex of degree 3 in every face; the graph is essentially 4-connected.

## **But not Hamiltonian**

Theorem [2016 Fabrici-Harant-Jendrol', 1992 Jackson-Wormald]: For infinitely many (maximal planar) essentially 4-connected polyhedral graphs G,  $circ(G) \leq 2(n+4)/3$ .

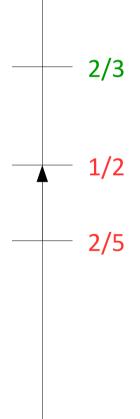


- As the red vertex set is independent, any cycle has at most twice the number of white vertices.
- As there are 2 = (n+4)/3-4 more red than white vertices, we miss 2 vertices.

# G polyhedral & essentially 4-connected

Theorem [2016 Fabrici-Harant-Jendrol']: circ(G)  $\geq$  (n+4)/2

true factor

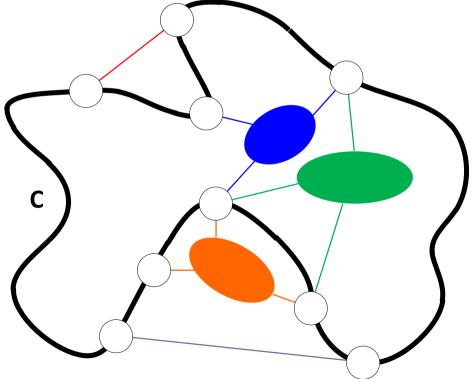


## **Alternative proof using 2-walks**

A cycle **C** of a plane 2-connected graph G is a **Tutte cycle** if

- every **C**-bridge has at most 3 attachments and
- every **C**-bridge containing an edge of the outer face has exactly 2 attachments.

An SDR of **C** is an SDR of the sets of attachment vertices of the **C**-bridges with 3 attachments.



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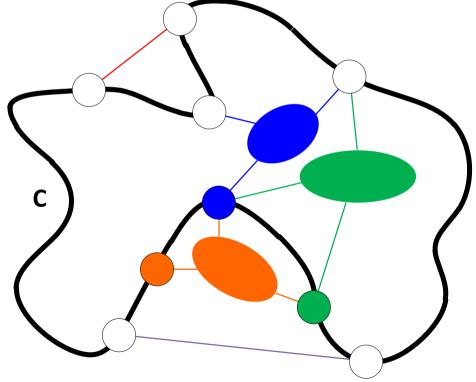
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An SDR of **C** is an SDR of the sets of attachment vertices of the **C**-bridges with 3 attachments.

Theorem [1995 Gao-Richter-Yu]: Every 3-connected plane graph has a 2-walk (+ a Tutte cycle **C** for which an SDR exists).

Is **C** long in essentially 4-connected graphs?

Not always!

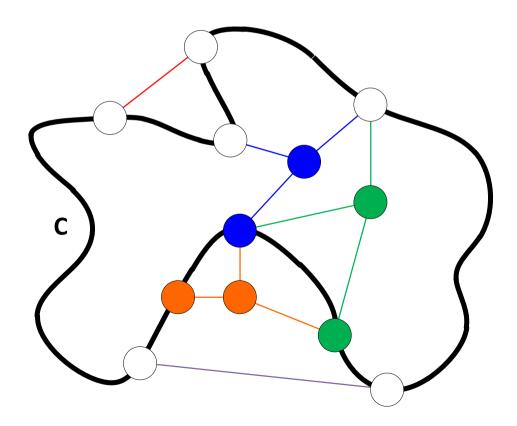


## **Alternative proof using 2-walks**

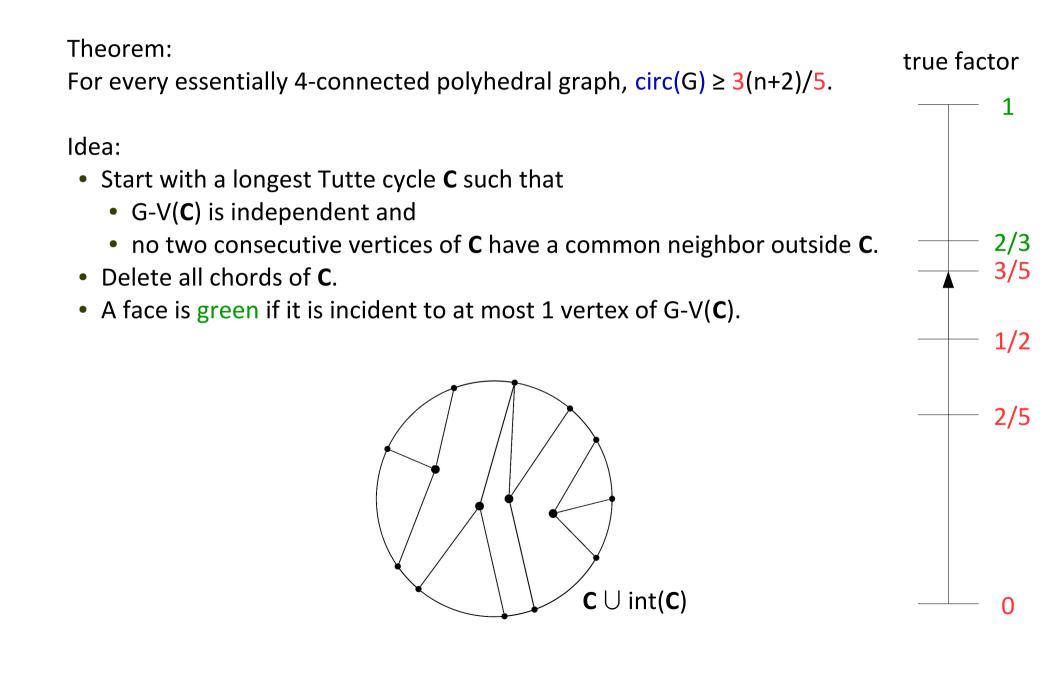
But there is a Tutte cycle C for which an SDR of C exists such that V(G)-V(C) is independent!

(Trick: Choose an appropriate Tutte cycle of length at least 4.)

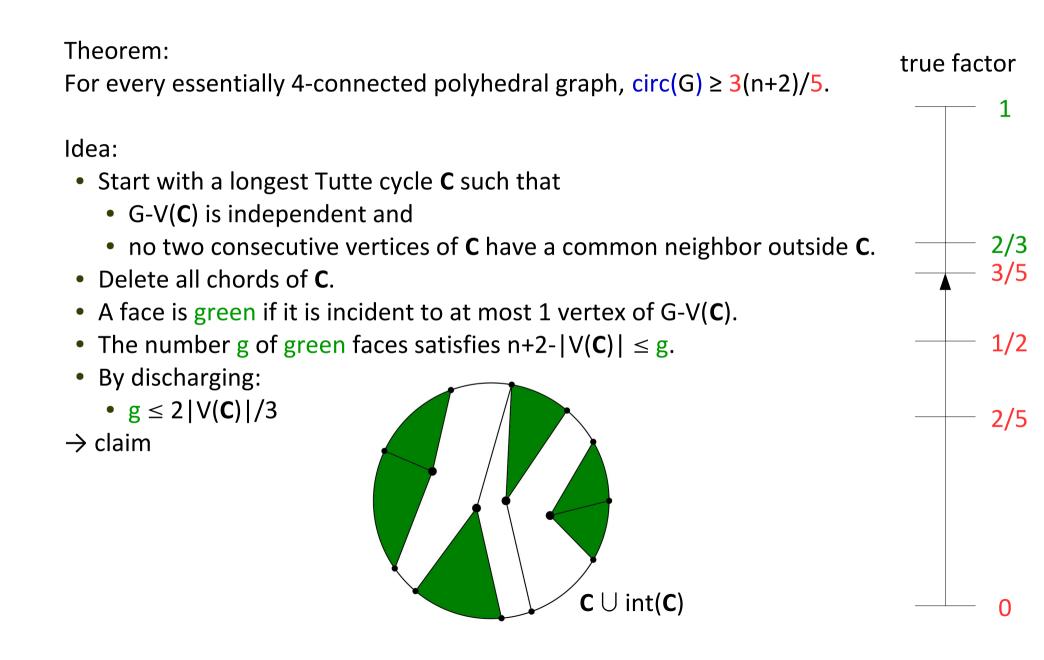
- $\rightarrow$  Every vertex not in C has degree 3.
- $\rightarrow$  |V(C)|  $\geq$  |V(G)-V(C)|
- $\rightarrow |V(C)| \ge n/2$



# **New Result**



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# Complexity

Theorem: For every essentially 4-connected polyhedral graph, a cycle of length at least 3(n+2)/5 can be computed in time  $O(n^2)$ .

Crucial parts of the previous proof: 
 instead, take a non-extendable one

- Start with a longest Tutte cycle **C** such that
- G-V(C) is independent and
  - no two consecutive vertices of **C** have a common neighbor outside **C**.

Theorem [2017 Schmid, S.]: A Tutte path of a 2-connected graph (in the general Sanders-variant) can be computed in time  $O(n^2)$ .

Thank you!