

Hamiltonicity of graphs on surfaces

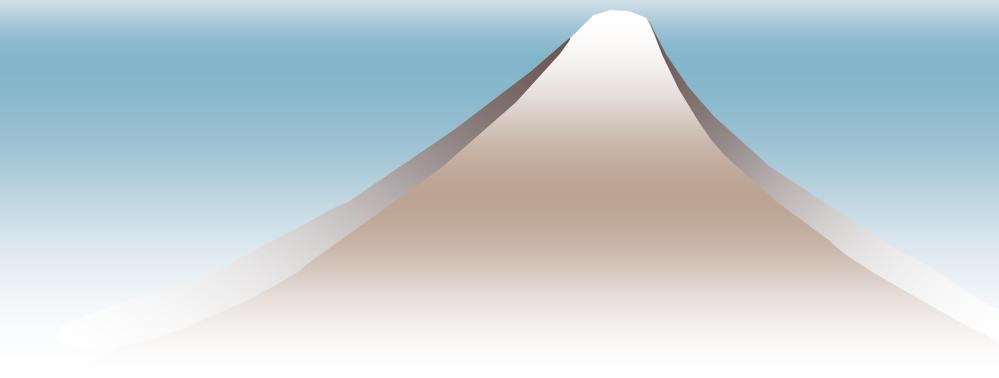
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(JST, ERATO, Kawarabayashi Large Graph Project)

Joint work with

Ken-ichi Kawarabayashi (National Institute of Informatics)



Hamiltonicity

G : Hamilton-connected

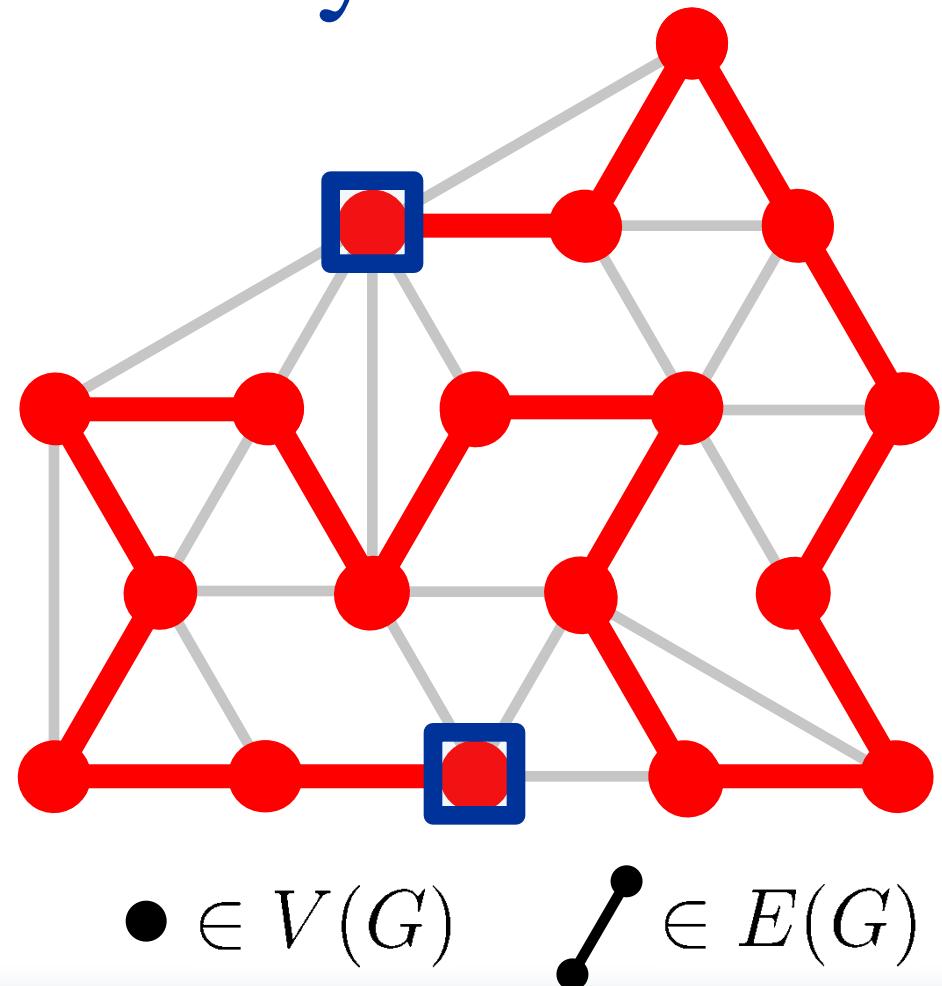


For \forall pair of vertices,
 \exists H-path between them

G : Hamilton-connected

$\Rightarrow \exists$ Hamilton cycle

$\Rightarrow \exists$ Hamilton path



Hamiltonicity

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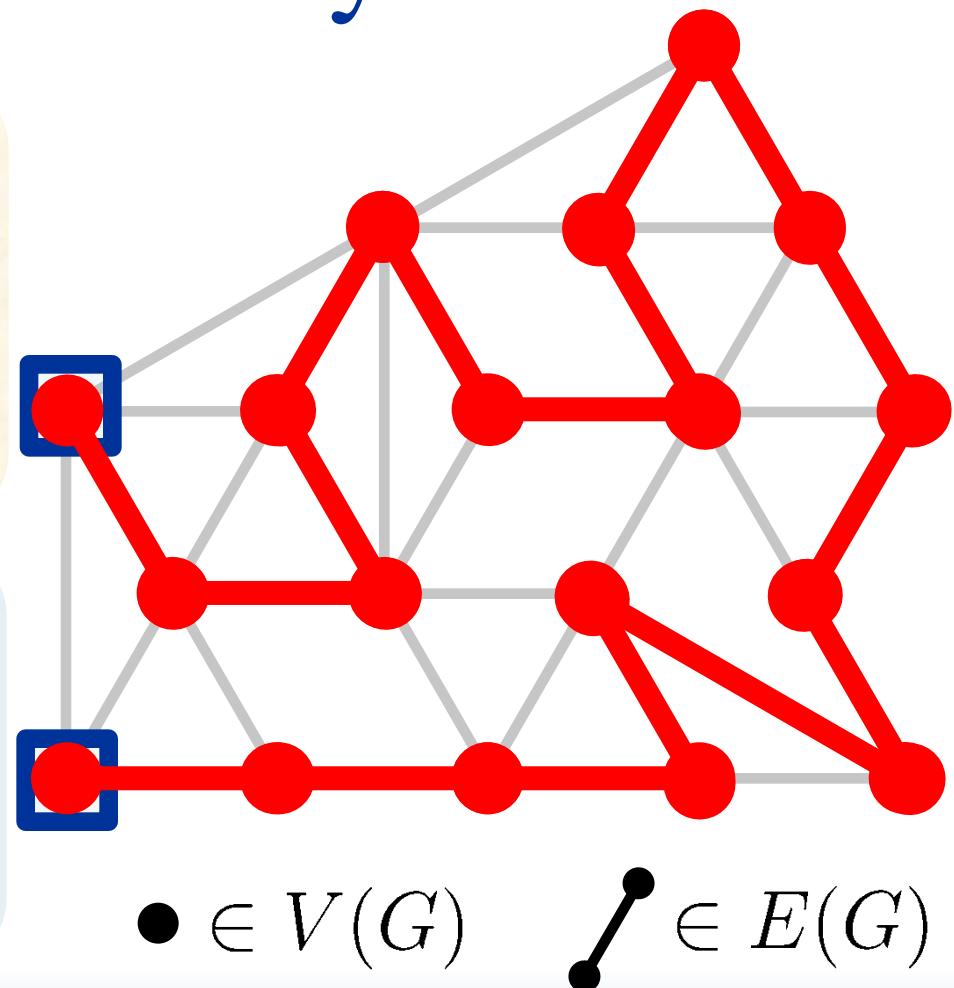


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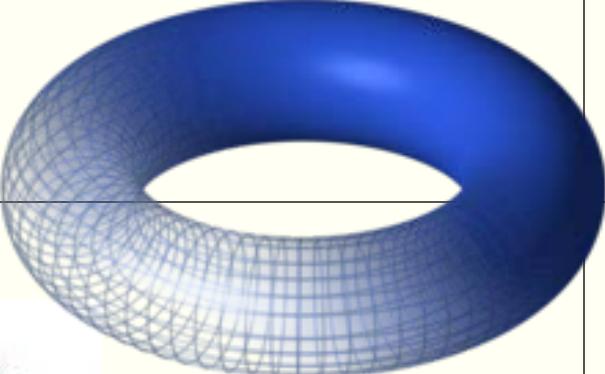
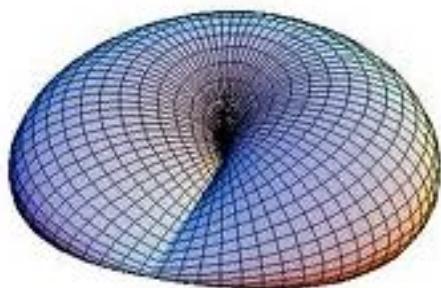
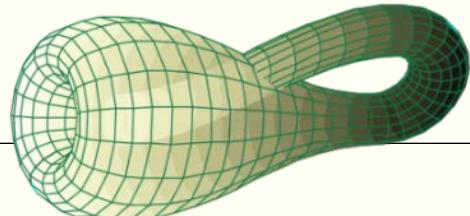
G : Hamilton-connected

$\Rightarrow \exists$ Hamilton cycle

$\Rightarrow \exists$ Hamilton path



Hamiltonicity of graphs on surfaces

	Plane $\chi = 2$	Proj. plane $\chi = 1$	Torus $\chi = 0$	K-bottle $\chi = 0$
\exists H-path				
\exists H-cycle				
H-conn				

Hamiltonicity of graphs on surfaces

4-connected			
	Plane $\chi = 2$	Proj. plane $\chi = 1$	Torus $\chi = 0$
\exists H-path	○	○	○ Thomas, Yu & Zang ('05)
\exists H-cycle	○ Tutte ('56)	○ Thomas & Yu ('94)	?
H-conn	○ Thomassen ('83)	○ K.K, Oz. ('14+)	✗

Tutte:

\forall 4-conn. **plane** graph
has a **H-cycle**

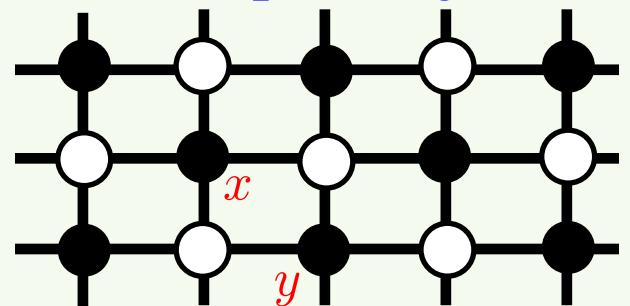
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Tutte:

\forall 4-conn. plane graph has a H-cycle

Take a bipartite grid.



\nexists H-path connecting x and y

Hamiltonicity of graphs on surfaces

	4-connected		5-connected	4-conn.	5-conn.
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\exists H-path	○	○	○ Thomas, Yu & Zang ('05)	○	?
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Some properties stronger than H-conn.

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Some properties stronger than H-conn.

C -Tutte path

T : C -Tutte path in G

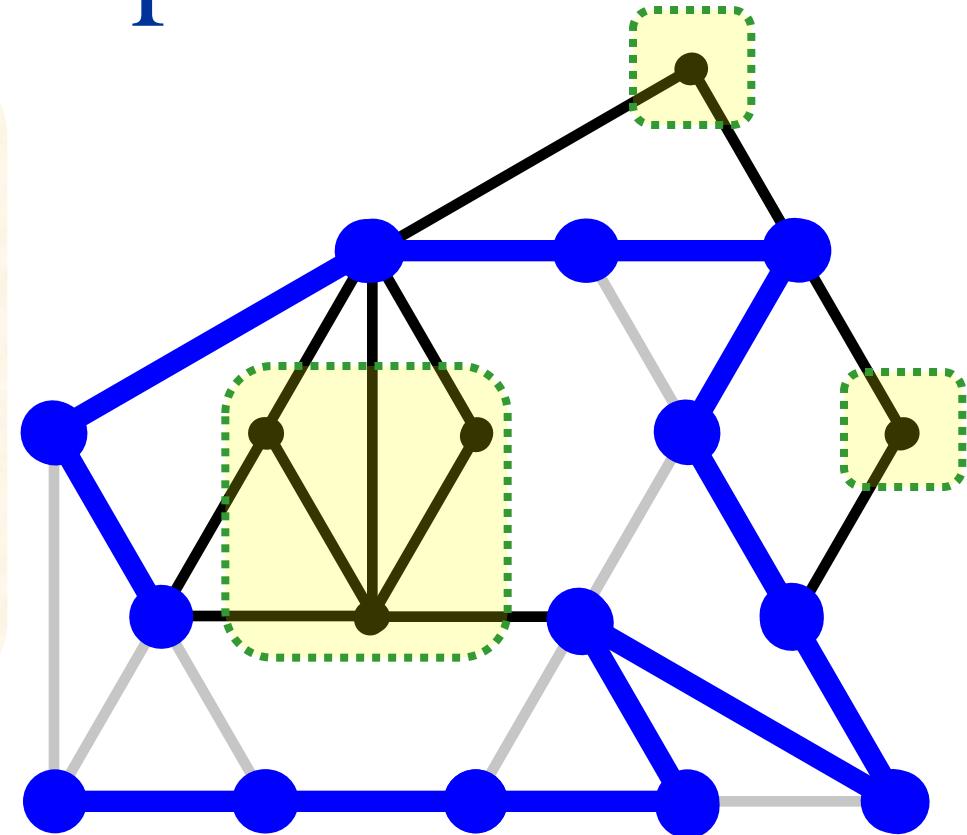
\Updownarrow

$\forall B$: compo. of $G - V(T)$

B has ≤ 3 neighbors on T

and ≤ 2 neighbors on T

if B contains a vertex in C



C -Tutte path in G

(C is the outer facial cycle)

C -Tutte path

T : C -Tutte path in G

\Updownarrow

$\forall B$: compo. of $G - V(T)$

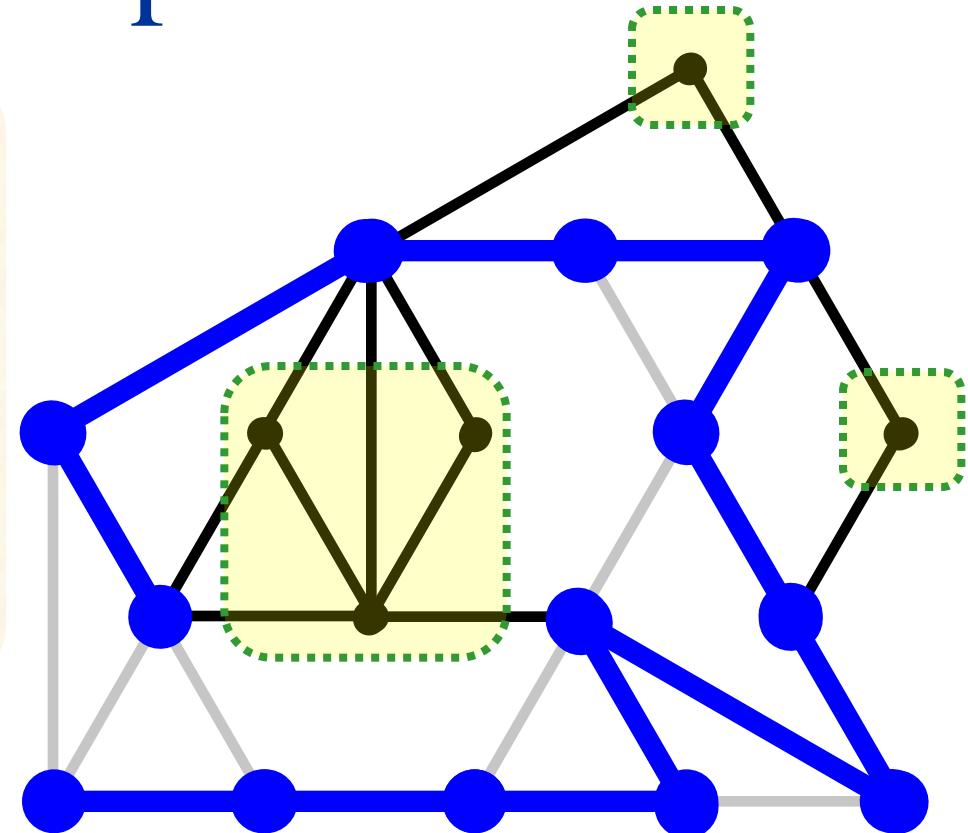
B has ≤ 3 neighbors on T

and ≤ 2 neighbors on T

if B contains a vertex in C

G : 4-conn. $|T| \geq 4$

$\Rightarrow T$: Hamiltonian path



C -Tutte path in G

(C is the outer facial cycle)

Idea of the proof

Condition $G : 2\text{-conn.}$

$x, y \in V(G)$ $C : \text{face containing } x$

Want to find :

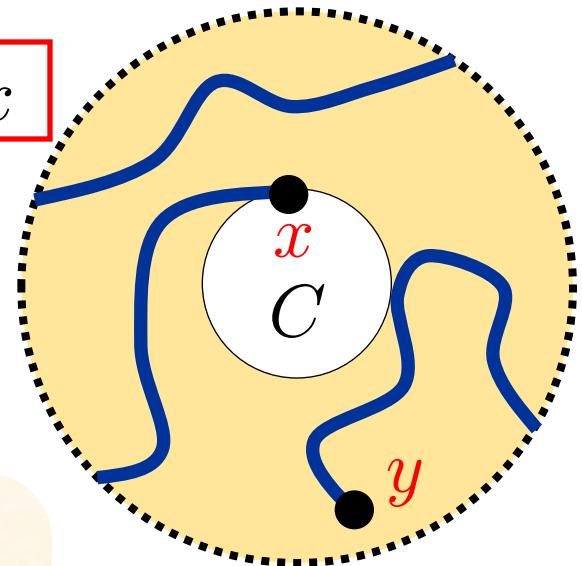
$T : C\text{-Tutte path between } x, y$

$T : C\text{-Tutte path}$

\Leftrightarrow For $\forall B : \text{compo. of } G - V(T),$

B has ≤ 3 neighbors on T

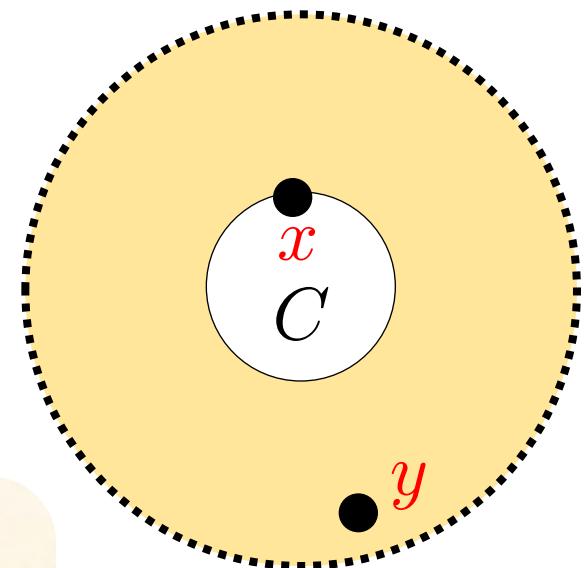
and ≤ 2 neighbors if B contains a vertex in C



Idea of the proof

Ordinary method

Use induction hypothesis to $G - V(C)$



T : C -Tutte path

\Leftrightarrow For $\forall B$: compo. of $G - V(T)$,

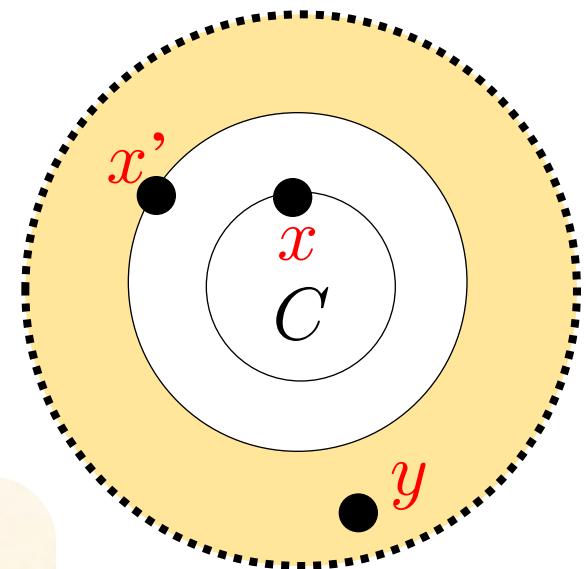
B has ≤ 3 neighbors on T

and ≤ 2 neighbors if B contains a vertex in C

Idea of the proof

Ordinary method

Use induction hypothesis to $G - V(C)$



T : C -Tutte path

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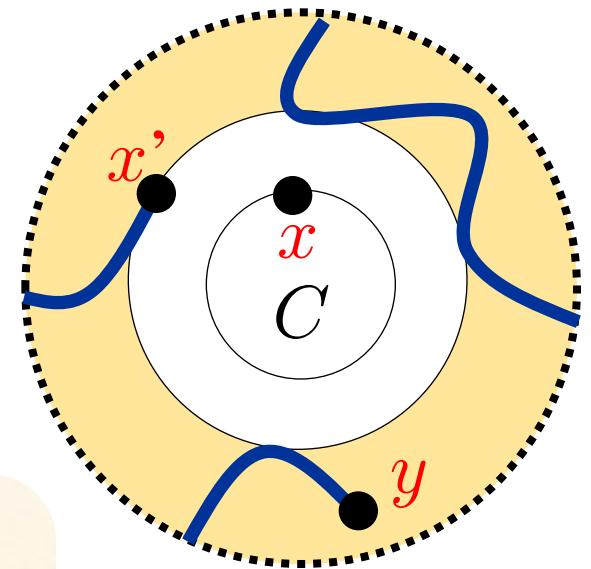
B has ≤ 3 neighbors on T

and ≤ 2 neighbors if B contains a vertex in C

Idea of the proof

Ordinary method

Use induction hypothesis to $G - V(C)$



T : C -Tutte path

\Leftrightarrow For $\forall B$: compo. of $G - V(T)$,

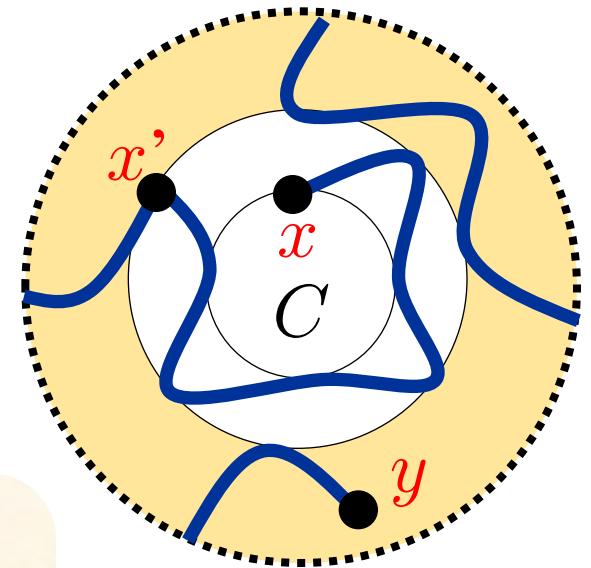
B has ≤ 3 neighbors on T

and ≤ 2 neighbors if B contains a vertex in C

Idea of the proof

Ordinary method

Use induction hypothesis to $G - V(C)$
and extend a C -Tutte path



T : C -Tutte path

\Leftrightarrow For $\forall B$: compo. of $G - V(T)$,

B has ≤ 3 neighbors on T

and ≤ 2 neighbors if B contains a vertex in C

Idea of the proof

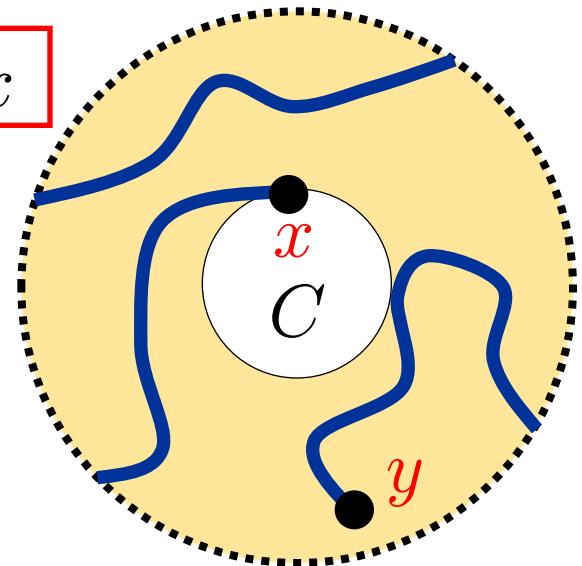
Condition G : 2-conn.

$x, y \in V(G)$ C : face containing x

Want to find :

T : C -Tutte path between x, y

Or



Idea of the proof

Condition G : 2-conn.

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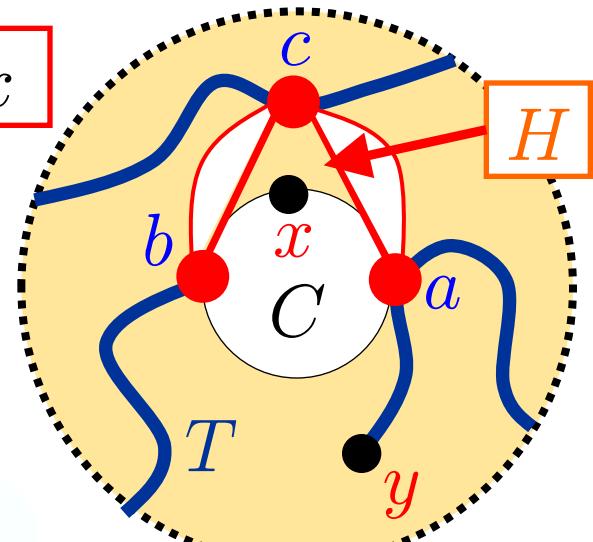
T : C -Tutte path between x, y

Or

$\exists \{a, b, c\}$: 3-cut of G separating x and y

$\exists H$: plane comp. of $G - \{a, b, c\}$ containing x (or $x = b$)

$\exists T$: C -Tutte path in $G - (H - \{a, b, c\})$ between b, y thr. a, c



Idea of the proof

Condition G : 2-conn.

$x, y \in V(G)$ | C : face containing x

Want to find :

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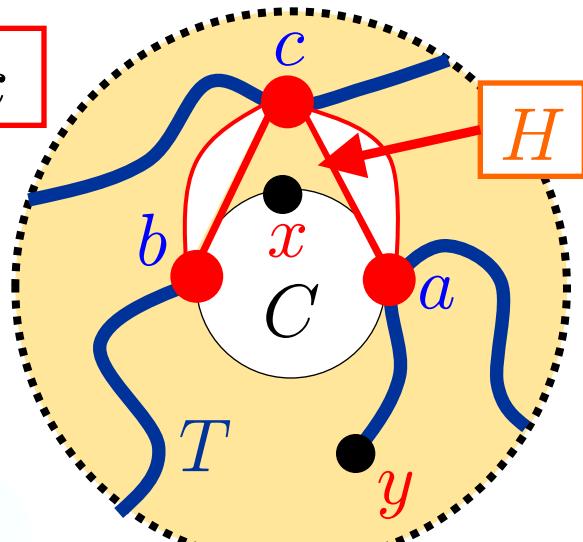
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$\exists \{a, b, c\} : \boxed{\text{3-cut}} \text{ of } G \text{ separating } x \text{ and } y$

$\exists H$: plane comp. of $G - \{a, b, c\}$ containing x (or $x = b$)

$\exists T$: **C-Tutte path** in $G - (H - \{a, b, c\})$ between b, y thr. a, c

We allow H : (unique) exceptional comp.

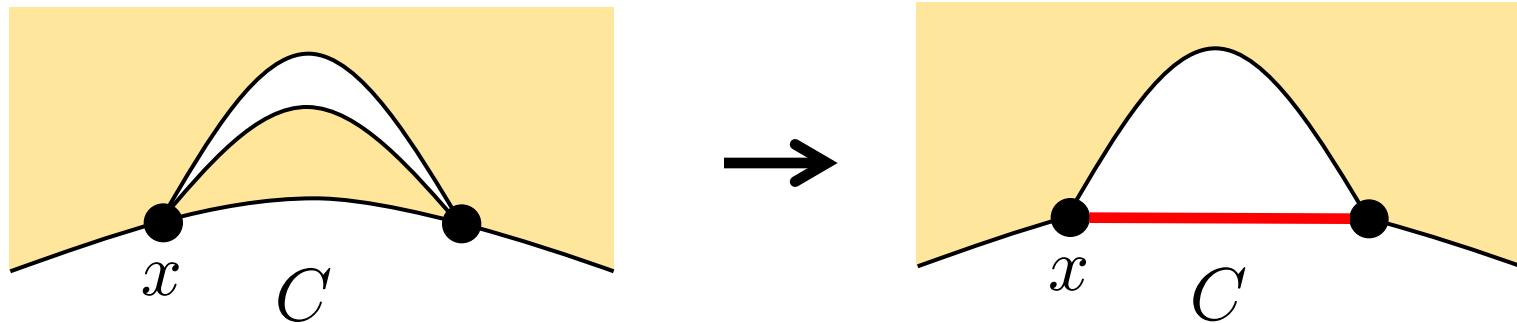


Idea of the proof

Induction on $|V(G)|$

Case I : $\exists S$: 2-cut with $x \in S$

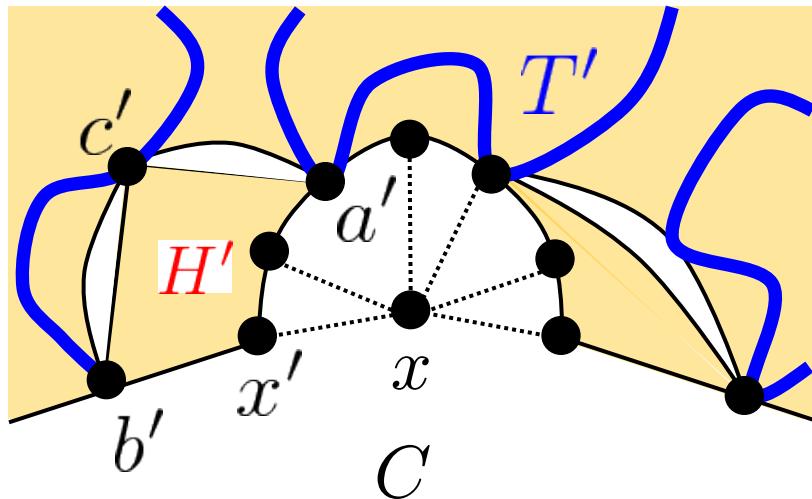
We use the reduction appearing in several situations.



Idea of the proof

Case II : $\nexists S$: 2-cut with $x \in S$

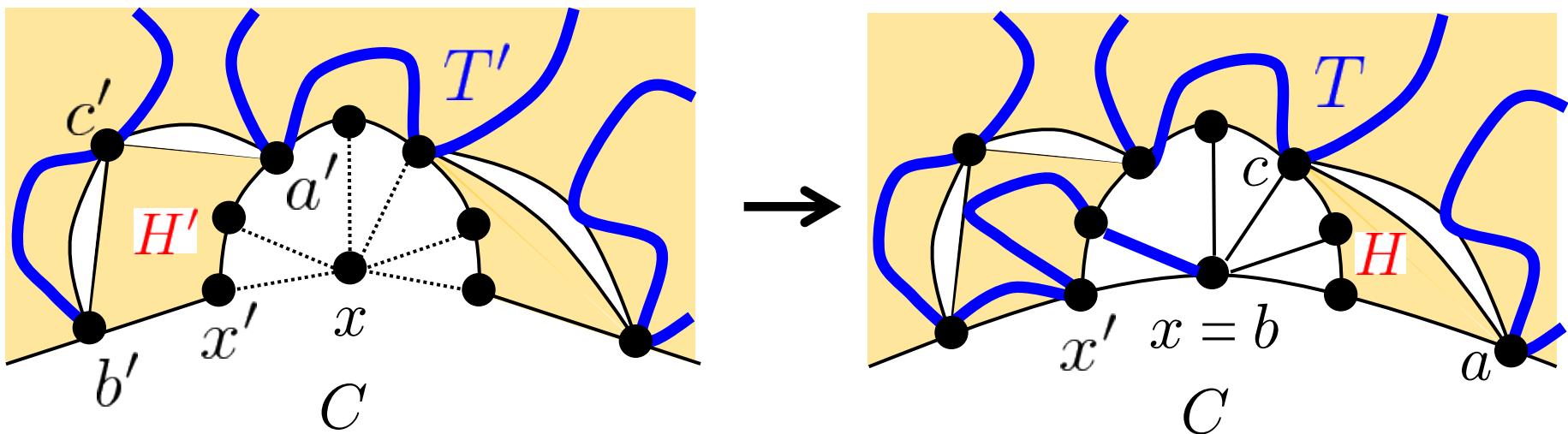
Then $G - x$ is also **2-connected**,
and use induction hypothesis to $G - x$



Idea of the proof

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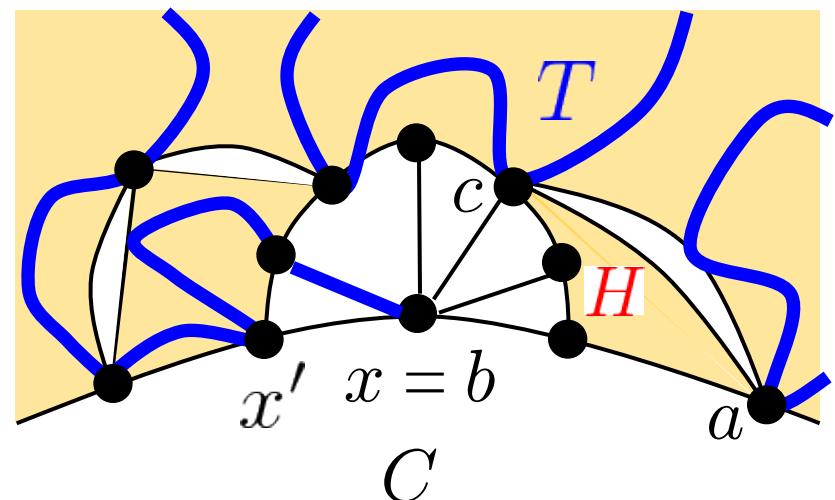
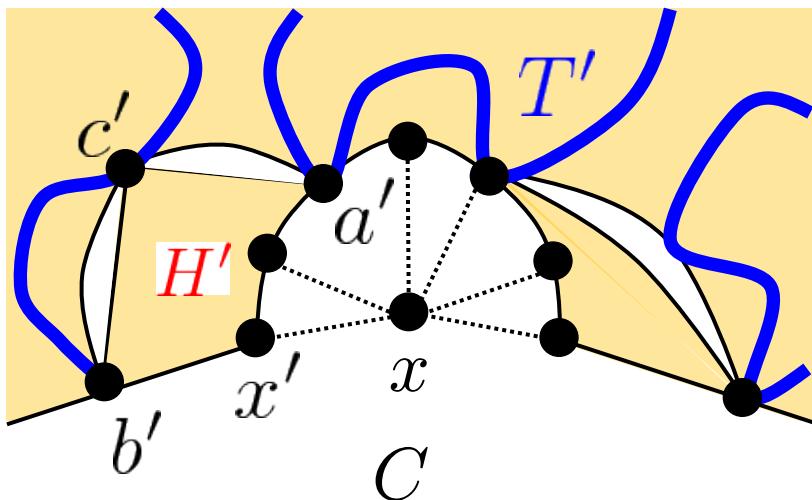
Idea of the proof

Case II : $\nexists S$: 2-cut with $x \in S$

Then $G - x$ is also **2-connected**,

and use induction hypothesis to $G - x$

Delete just one vertex



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Thank you for your attention

